## Context-free Grammars: In-class Exercise

(1) Consider the CFG $G$ with $S^{\prime}$ as the start symbol:

$$
\begin{aligned}
S^{\prime} & \rightarrow S \mid \epsilon \\
S & \rightarrow T \mid(N, C) \\
C & \rightarrow C, S \mid S \\
T & \rightarrow a|b| c \\
N & \rightarrow x|y| z
\end{aligned}
$$

a. List the set of terminal symbols and the set of non-terminal symbols in $G$.
b. For each of the following strings, write down true if the string is in the language $L(G)$ generated by $G$, false otherwise.

1. y
2. c
3. $(x)$
4. $(x, y)$
5. $(z, a, b, a, b, c)$
6. $(x, a,(y, b), c)$
7. $(x,(y, a),(z, b))$
8. ( $x,(x,(x,(x, a))$
(2) One of the rules in the CFG below is redundant: any sentence that can be generated using this rule can already be generated by a combination of other rules. Write down the redundant rule.

| $\mathrm{S} \rightarrow \mathrm{NPVP}$ | IV $\rightarrow$ runs | $\mathrm{N} \rightarrow$ John |
| :---: | :---: | :---: |
| $\mathrm{NP} \rightarrow \mathrm{N}$ | IV $\rightarrow$ sits | $\mathrm{N} \rightarrow$ he |
| $\mathrm{NP} \rightarrow \mathrm{DN}$ | TV $\rightarrow$ chases | $\mathrm{N} \rightarrow$ Mary |
| $\mathrm{VP} \rightarrow$ VP PP | TV $\rightarrow$ eats | $\mathrm{N} \rightarrow$ dog |
| VP $\rightarrow$ VP CONJ VP | TV $\rightarrow$ catches | $\mathrm{N} \rightarrow$ tree |
| $\mathrm{VP} \rightarrow$ IV | TV $\rightarrow$ tells | $\mathrm{N} \rightarrow$ squirrel |
| VP $\rightarrow$ IV PP | TV $\rightarrow$ sees | $\mathrm{D} \rightarrow$ the |
| VP $\rightarrow$ TV NP | CONJ $\rightarrow$ and |  |
| VP $\rightarrow$ TVCS | $\mathrm{C} \rightarrow$ that |  |
| NP $\rightarrow$ NP CONJ NP | $\mathrm{P} \rightarrow$ in |  |
| $\mathrm{PP} \rightarrow \mathrm{P}$ | $\mathrm{P} \rightarrow$ away |  |
| $\mathrm{PP} \rightarrow \mathrm{PNP}$ |  |  |

(3) Consider the family of CFGs $G_{k}$ with $S$ as the start symbol and $k$ is some arbitrary non-zero positive integer such that $G_{1}, G_{2}, G_{3}, \ldots$ are individual CFGs with the rules:

$$
\begin{aligned}
& S \rightarrow A B \\
& B \rightarrow C A A \\
& C \rightarrow c \\
& A \rightarrow a_{i} \text { defines } i \text { rules, where } i \in[1, k]
\end{aligned}
$$

For example, in $G_{3}$ the rules with $A$ as left-hand side are: $A \rightarrow a_{1}\left|a_{2}\right| a_{3}$ with three terminal symbols.
a. Provide the number of terminal symbols in a grammar $G_{k}$.
b. If the string $a_{4} \mathrm{Ca}_{3} a_{2}$ is accepted by grammar $G_{3}$ then provide a derivation for it.
c. If the string $a_{4} \mathrm{ca}_{3} a_{2}$ is accepted by grammar $G_{4}$ then provide a derivation for it.
d. Provide the total number of strings that can be generated for a grammar $G_{k}$.
(4) Consider a treebank which consists of three tree types: $T_{1}, T_{2}, T_{3}$. In this treebank these tree types are repeated multiple times. By counting the number of times each tree type was observed, we discover that each tree type occurs with the following probability:

$$
\begin{array}{ll}
p_{1} & T_{1}=(\mathrm{S}(\mathrm{~B} a)(\mathrm{C} a \mathrm{a})) \\
p_{2} & T_{2}=(\mathrm{S}(\mathrm{~B} a \mathrm{a})) \\
p_{3} & T_{3}=(\mathrm{S}(\mathrm{C} a \mathrm{a}))
\end{array}
$$

a. From the treebank shown above, extract a probabilistic CFG (PCFG) $G$.

Assume that the rule $\mathrm{S} \rightarrow \mathrm{BC}$ appears in $G$ with probability $p_{1}$ and $p_{1}+p_{2}+p_{3}=1$.
b. Provide the tree set $\mathcal{T}$ for the $\mathrm{CFG} G$.
c. Provide the language $\mathcal{L}$ (the set of strings) for the $\mathrm{CFG} G$.
d. Let $p_{1}=0.2, p_{2}=0.1, p_{3}=0.7$. Find the parse tree with highest probability according to PCFG $G$ for the input string $a a$. Note that tree $T_{2}$ in the treebank is a tree that has yield $a a$. Write down if the tree you find is the same as $T_{2}$.

