# Natural Language Processing 

Anoop Sarkar<br>anoopsarkar.github.io/nlp-class

Simon Fraser University

September 5, 2019

# Natural Language Processing 

Anoop Sarkar<br>anoopsarkar.github.io/nlp-class

Simon Fraser University

Part 1: Ambiguity

## Context Free Grammars and Ambiguity

$$
\begin{array}{clc}
S & \rightarrow & N P V P \\
V P & \rightarrow & V N P \\
V P & \rightarrow & V P P P \\
P P & \rightarrow P N P \\
N P & \rightarrow & N P P P \\
N P & \rightarrow & \text { Calvin } \\
N P & \rightarrow & \text { monsters } \\
N P & \rightarrow & \text { school } \\
V & \rightarrow & \text { imagined } \\
P & \rightarrow & \text { in }
\end{array}
$$

What is the analysis using the above grammar for:
Calvin imagined monsters in school

## Context Free Grammars and Ambiguity

Calvin imagined monsters in school

```
(S (NP Calvin)
    (VP (V imagined)
            (NP (NP monsters)
            (PP (P in)
                (NP school)))))
```

(S (NP Calvin)
(VP (VP (V imagined)
(NP monsters))
(PP (P in)
(NP school))))

Which one is more plausible?

## Context Free Grammars and Ambiguity

Calvin imagined monsters in school


Calvin imagined monsters in school


## Ambiguity Kills (your parser)

```
natural language learning course
(run demos/parsing-ambiguity.py)
((natural language) (learning course))
(((natural language) learning) course)
((natural (language learning)) course)
(natural (language (learning course)))
(natural ((language learning) course))
```

- Some difficult issues:
- Which one is more plausible?
- How many analyses for a given input?
- Computational complexity of parsing language


## Number of derivations

CFG rules $\{\mathrm{N} \rightarrow \mathrm{NN}, \mathrm{N} \rightarrow \mathrm{a}\}$

| $n: a^{n}$ | number of parses |
| :---: | :---: |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 5 |
| 5 | 14 |
| 6 | 42 |
| 7 | 132 |
| 8 | 429 |
| 9 | 1430 |
| 10 | 4862 |
| 11 | 16796 |

## CFG Ambiguity

- Number of parses in previous table is an integer series, known as the Catalan numbers
- Catalan numbers have a closed form:

$$
\operatorname{Cat}(n)=\frac{1}{n+1}\binom{2 n}{n}
$$

- $\binom{a}{b}$ is the binomial coefficient

$$
\binom{a}{b}=\frac{a!}{(b!(a-b)!)}
$$

## Catalan numbers

- Why Catalan numbers? Cat $(\mathrm{n})$ is the number of ways to parenthesize an expression of length $n$ with two conditions:

1. there must be equal numbers of open and close parens
2. they must be properly nested so that an open precedes a close

- ( $(\mathrm{ab}) \mathrm{c}) \mathrm{d}(\mathrm{a}(\mathrm{bc})) \mathrm{d}(\mathrm{ab})(\mathrm{cd}) \mathrm{a}((\mathrm{bc}) \mathrm{d}) \mathrm{a}(\mathrm{b}(\mathrm{cd}))$
- For an expression of with $n$ ways to form constituents there are a total of $2 n$ choose $n$ parenthesis pairs. Then divide by $n+1$ to remove invalid parenthesis pairs.
- For more details see (Church and Patil, CL Journal, 1982)


# Natural Language Processing 

Anoop Sarkar<br>anoopsarkar.github.io/nlp-class

Simon Fraser University

## Part 2: Context Free Grammars

## Context-Free Grammars

- A CFG is a 4-tuple: $(N, T, R, S)$, where
- $N$ is a set of non-terminal symbols,
- $T$ is a set of terminal symbols which can include the empty string $\epsilon$. $T$ is analogous to $\Sigma$ the alphabet in FSAs.
- $R$ is a set of rules of the form $A \rightarrow \alpha$, where $A \in N$ and $\alpha \in\{N \cup T\}^{*}$
- $S$ is a set of start symbols, $S \in N$


## Context-Free Grammars

- Here's an example of a CFG, let's call this one $G$ :

1. $S \rightarrow a S b$
2. $S \rightarrow \epsilon$

- What is the language of this grammar, which we will call $L(G)$, the set of strings generated by this grammar How? Notice that there cannot be any FSA that corresponds exactly to this set of strings $L(G)$ Why?
- What is the tree set or derivations produced by this grammar?


## Context-Free Grammars

- This notion of generating both the strings and the trees is an important one for Computational Linguistics
- Consider the trees for the grammar $G^{\prime}$ :
$P=\{S \rightarrow A A, A \rightarrow a A, A \rightarrow A b, A \rightarrow \epsilon\}$,
$\Sigma=\{a, b\}, N=\{S, A\}, T=\{a, b, \epsilon\}, S=\{S\}$
- Why is it called context-free grammar?


## Context-Free Grammars

- Can the grammar $G^{\prime}$ produce only trees with equal height subtrees on the left and right?



## Parse Trees

Consider the grammar with rules:

$$
\begin{aligned}
S & \rightarrow N P V P \\
N P & \rightarrow P R P \\
N P & \rightarrow D T N P B \\
V P & \rightarrow V B P N P \\
N P B & \rightarrow N N N N \\
P R P & \rightarrow I \\
V B P & \rightarrow \text { prefer } \\
D T & \rightarrow a \\
N N & \rightarrow \text { morning } \\
N N & \rightarrow \text { flight }
\end{aligned}
$$

## Parse Trees



## Parse Trees: Equivalent Representations

- (S (NP (PRP I)) (VP (VBP prefer) (NP (DT a) (NPB (NN morning) (NN flight)))))
- [s [NP [PRP I ] ] [VP [VBP prefer ] [NP [DT a ] [NPB [NN morning ] [NN flight ] ] ] ] ]


## Ambiguous Grammars

- $S \rightarrow S S$
- $S \rightarrow a$
- Given the above rules, consider the input aaa, what are the valid parse trees?
- Now consider the input aaaa


## Inherently Ambiguous Languages

- Consider the following context-free grammar:
- $S \rightarrow S 1 \mid S 2$
- $S 1 \rightarrow a X d \mid \epsilon$
- $X \rightarrow b X c \mid \epsilon$
- $S 2 \rightarrow Y Z \mid \epsilon$
- $Y \rightarrow a Y b \mid \epsilon$
- $Z \rightarrow c Z d \mid \epsilon$
- Now parse the input string abcd with this grammar
- Notice that we get two parse trees (one with the S1 sub-grammar and another with the $S 2$ subgrammar).


# Natural Language Processing 

Anoop Sarkar<br>anoopsarkar.github.io/nlp-class

Simon Fraser University

Part 3: Structural Ambiguity

## Ambiguity

- Part of Speech ambiguity
saw $\rightarrow$ noun
saw $\rightarrow$ verb
- Structural ambiguity: Prepositional Phrases

I saw (the man) with the telescope
I saw (the man with the telescope)

- Structural ambiguity: Coordination
a program to promote safety in ((trucks) and (minivans))
a program to promote ((safety in trucks) and (minivans))
((a program to promote safety in trucks) and (minivans))


## Ambiguity $\leftarrow$ attachment choice in alternative parses



## Ambiguity in Prepositional Phrases

- noun attach: I bought the shirt with pockets
- verb attach: I washed the shirt with soap
- As in the case of other attachment decisions in parsing: it depends on the meaning of the entire sentence - needs world knowledge, etc.
- Maybe there is a simpler solution: we can attempt to solve it using heuristics or associations between words


## Structure Based Ambiguity Resolution

- Right association: a constituent (NP or PP) tends to attach to another constituent immediately to its right (Kimball 1973)
- Minimal attachment: a constituent tends to attach to an existing non-terminal using the fewest additional syntactic nodes (Frazier 1978)
- These two principles make opposite predictions for prepositional phrase attachment
- Consider the grammar:

$$
\begin{align*}
& V P \rightarrow V N P P P  \tag{1}\\
& N P \rightarrow N P P P \tag{2}
\end{align*}
$$

for input: I [VP saw [NP the man ... [PP with the telescope ], RA predicts that the PP attaches to the NP, i.e. use rule (2), and MA predicts V attachment, i.e. use rule (1)

## Structure Based Ambiguity Resolution

- Garden-paths look structural: The emergency crews hate most is domestic violence
- Neither MA or RA account for more than $55 \%$ of the cases in real text
- Psycholinguistic experiments using eyetracking show that humans resolve ambiguities as soon as possible in the left to right sequence using the words to disambiguate
- Garden-paths are caused by a combination of lexical and structural effects:
The flowers delivered for the patient arrived


## Ambiguity Resolution: Prepositional Phrases in English

- Learning Prepositional Phrase Attachment: Annotated Data

| v | n 1 | p | n 2 | Attachment |
| :---: | :---: | :---: | :---: | :---: |
| join | board | as | director | V |
| is | chairman | of | N.V. | N |
| using | crocidolite | in | filters | V |
| bring | attention | to | problem | V |
| is | asbestos | in | products | N |
| making | paper | for | filters | N |
| including | three | with | cancer | N |
| $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ | $\vdots$ |

## Prepositional Phrase Attachment

| Method | Accuracy |
| :--- | :---: |
| Always noun attachment | 59.0 |
| Most likely for each preposition | 72.2 |
| Average Human (4 head words only) | 88.2 |
| Average Human (whole sentence) | 93.2 |

## Some other studies

- Toutanova, Manning, and Ng, 2004: 87.54\% using some external knowledge (word classes)
- Merlo, Crocker and Berthouzoz, 1997: test on multiple PPs
- generalize disambiguation of 1 PP to 2-3 PPs
- 14 structures possible for 3PPs assuming a single verb
- all 14 are attested in the Penn WSJ Treebank
- 1PP: 84.3\% 2PP: 69.6\% 3PP: 43.6\%
- Belinkov+ TACL 2014: Neural networks for PP attachment (multiple candidate heads)
- NN model (no extra data): 86.6\%
- NN model (lots of raw data for word vectors): $88.7 \%$
- NN model with parser and lots of raw data: $90.1 \%$
- This experiment is still only part of the real problem faced in parsing English. Plus other sources of ambiguity in other languages


# Natural Language Processing 

Anoop Sarkar<br>anoopsarkar.github.io/nlp-class

Simon Fraser University

Part 4: Weighted Context Free Grammars

## Treebanks

- What is the CFG that can be extracted from this single tree:
(S (NP (Det the) (NP man))
(VP (VP (V played)
(NP (Det a) (NP game)))
(PP (P with)
(NP (Det the) (NP dog)))))


## PCFG

| $S$ | $\rightarrow$ | $N P V P$ | $c=1$ |
| :---: | :--- | :---: | :---: |
| $N P$ | $\rightarrow$ | $\operatorname{Det~} N P$ | $c=3$ |
| $N P$ | $\rightarrow$ | man | $c=1$ |
| $N P$ | $\rightarrow$ | game | $c=1$ |
| $N P$ | $\rightarrow$ | dog | $c=1$ |
| $V P$ | $\rightarrow$ | $V P P P$ | $c=1$ |
| $V P$ | $\rightarrow$ | $V N P$ | $c=1$ |
| $P P$ | $\rightarrow$ | $P N P$ | $c=1$ |
| $D e t$ | $\rightarrow$ | the | $c=2$ |
| Det | $\rightarrow$ | $a$ | $c=1$ |
| $V$ | $\rightarrow$ | played | $c=1$ |
| $P$ | $\rightarrow$ | with | $c=1$ |

- We can do this with multiple trees. Simply count occurrences of CFG rules over all the trees.
- A repository of such trees labelled by a human is called a TreeBank.


## Probabilistic CFG (PCFG)

$$
\begin{array}{clcc}
S & \rightarrow & N P V P & 1 \\
V P & \rightarrow & V N P & 0.9 \\
V P & \rightarrow & V P P P & 0.1 \\
P P & \rightarrow & P N P & 1 \\
N P & \rightarrow & N P P P & 0.25 \\
N P & \rightarrow & \text { Calvin } & 0.25 \\
N P & \rightarrow & \text { monsters } & 0.25 \\
N P & \rightarrow & \text { school } & 0.25 \\
V & \rightarrow & \text { imagined } & 1 \\
P & \rightarrow & \text { in } & 1
\end{array}
$$

$$
P(\text { input })=\sum_{\text {tree }} P(\text { tree } \mid \text { input })
$$

$P($ Calvin imagined monsters in school $)=$ ?
Notice that $P(V P \rightarrow V N P)+P(V P \rightarrow V P P P)=1.0$

## Probabilistic CFG (PCFG)

$P($ Calvin imagined monsters in school $)=$ ?

(S (NP Calvin)
(VP (VP (V imagined)
(NP monsters))
(PP (P in)
(NP school))))

## Probabilistic CFG (PCFG)

```
(S (NP Calvin)
    (VP (V imagined)
        (NP (NP monsters)
        (PP (P in)
                                (NP school)))))
```

$$
\begin{aligned}
P\left(\text { tree }_{1}\right)= & P(S \rightarrow N P V P) \times P(N P \rightarrow \text { Calvin }) \times P(V P \rightarrow V N P) \times \\
& P(V \rightarrow \text { imagined }) \times P(N P \rightarrow N P P P) \times P(N P \rightarrow \text { monsters }) \times \\
& P(P P \rightarrow P N P) \times P(P \rightarrow \text { in }) \times P(N P \rightarrow \text { school }) \\
= & 1 \times 0.25 \times 0.9 \times 1 \times 0.25 \times 0.25 \times 1 \times 1 \times 0.25=.003515625
\end{aligned}
$$

## Probabilistic CFG (PCFG)

```
(S (NP Calvin)
    (VP (VP (V imagined)
        (NP monsters))
        (PP (P in)
        (NP school))))
\[
\begin{aligned}
P\left(\text { tree }_{2}\right)= & P(S \rightarrow N P V P) \times P(N P \rightarrow \text { Calvin }) \times P(V P \rightarrow V P P P) \times \\
& P(V P \rightarrow V N P) \times P(V \rightarrow \text { imagined }) \times P(N P \rightarrow \text { monsters }) \times \\
& P(P P \rightarrow P N P) \times P(P \rightarrow \text { in }) \times P(N P \rightarrow \text { school }) \\
= & 1 \times 0.25 \times 0.1 \times 0.9 \times 1 \times 0.25 \times 1 \times 1 \times 0.25=.00140625
\end{aligned}
\]
```


## Probabilistic CFG (PCFG)

$P($ Calvin imagined monsters in school $)=P\left(\right.$ tree $\left._{1}\right)+P\left(\right.$ tree $\left._{2}\right)$

$$
\begin{aligned}
& =.003515625+.00140625 \\
& =.004921875
\end{aligned}
$$

$$
\text { Most likely tree is tree } 1=\underset{\text { tree }}{\arg \max } P(\text { tree } \mid \text { input })
$$

(S (NP Calvin)
(VP (V imagined)
(NP (NP monsters)
(PP (P in)
(NP school)))))
(S (NP Calvin)
(VP (VP (V imagined)
(NP monsters))
(PP (P in)
(NP school))))

## Probabilistic Context-Free Grammars (PCFG)

- A PCFG is a 4-tuple: $(N, T, R, S)$, where
- $N$ is a set of non-terminal symbols,
- $T$ is a set of terminal symbols which can include the empty string $\epsilon$. $T$ is analogous to $\Sigma$ the alphabet in FSAs.
- $R$ is a set of rules of the form $A \rightarrow \alpha$, where $A \in N$ and $\alpha \in\{N \cup T\}^{*}$
- $P(R)$ is the probability of rule $R: A \rightarrow \alpha$ such that $\sum_{\alpha} P(A \rightarrow \alpha)=1.0$
- $S$ is a set of start symbols, $S \in N$


## PCFG

- Central condition: $\sum_{\alpha} P(A \rightarrow \alpha)=1$
- Called a proper PCFG if this condition holds
- Note that this means $P(A \rightarrow \alpha)=P(\alpha \mid A)=\frac{f(A, \alpha)}{f(A)}$
- $P(T \mid S)=\frac{P(T, S)}{P(S)}=P(T, S)=\prod_{i} P\left(R H S_{i} \mid L H S_{i}\right)$


## PCFG

- What is the PCFG that can be extracted from this single tree:
(S (NP (Det the) (NP man))
(VP (VP (V played)
(NP (Det a) (NP game)))
(PP (P with)
(NP (Det the) (NP dog)))))
- How many different rhs $\alpha$ exist for $A \rightarrow \alpha$ where $A$ can be $S$, NP, VP, PP, Det, N, V, P


## PCFG

| $S$ | $\rightarrow$ | $N P V P$ | $c=1$ | $p=1 / 1=1.0$ |
| :---: | :---: | :---: | :---: | :---: | :--- |
| $N P$ | $\rightarrow$ | $\operatorname{Det~NP}$ | $c=3$ | $p=3 / 6=0.5$ |
| $N P$ | $\rightarrow$ | man | $c=1$ | $p=1 / 6=0.1667$ |
| $N P$ | $\rightarrow$ | game | $c=1$ | $p=1 / 6=0.1667$ |
| $N P$ | $\rightarrow$ | $\operatorname{dog}$ | $c=1$ | $p=1 / 6=0.1667$ |
| $V P$ | $\rightarrow$ | $V P P P$ | $c=1$ | $p=1 / 2=0.5$ |
| $V P$ | $\rightarrow$ | $V N P$ | $c=1$ | $p=1 / 2=0.5$ |
| $P P$ | $\rightarrow$ | $P N P$ | $c=1$ | $p=1 / 1=1.0$ |
| Det | $\rightarrow$ | the | $c=2$ | $p=2 / 3=0.67$ |
| Det | $\rightarrow$ | $a$ | $c=1$ | $p=1 / 3=0.33$ |
| $V$ | $\rightarrow$ | played | $c=1$ | $p=1 / 1=1.0$ |
| $P$ | $\rightarrow$ | with | $c=1$ | $p=1 / 1=1.0$ |

- We can do this with multiple trees. Simply count occurrences of CFG rules over all the trees.
- A repository of such trees labelled by a human is called a TreeBank.

